

# COVERING SIMPLY CONNECTED REGIONS BY RECTANGLES

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We prove that the ratio of the minimum number of rectangles covering a simply connected board (polyomino)  $B$  and the maximum number of points in  $B$  no two of which are contained in a common rectangle is less than 2.

A *board* (polyomino) is a finite set of unit squares lying in the plane whose corners have integer coordinates. In other words, a board is a lattice polygon with vertical and horizontal sides. A *rectangle* is a subset of the board whose union is rectangular. A *cover* of a board  $B$  is a collection of rectangles whose union is equal to  $B$ . (That is, the rectangles of a cover may overlap but they must be contained in  $B$ .) An *antirectangle* in  $B$  is a set of squares in  $B$  no two of which are contained in a common rectangle. It is obvious that any cover has to contain at least as many rectangles as any antirectangle has squares. Thus if  $\theta(B)$  is the minimum number of rectangles in a cover of  $B$  and  $\alpha(B)$  is the maximum number of squares in an antirectangle of  $B$  then  $\alpha(B) \leq \theta(B)$ . V. Chvátal originally conjectured that  $\alpha(B) = \theta(B)$  holds for any finite board  $B$ . In general, this is false. First E. Szemerédi [4] found a counterexample with a "hole" (Figure 2), then F.R.K. Chung [2] found the simply connected counterexample in Figure 1. Then S. Chaiken, D. J. Kleitman, M. Saks and J. Shearer [1] proved a weakened version of the conjecture that equality does hold if  $B$  is horizontally and vertically convex i.e. whenever two squares in  $B$  are on the same horizontal or vertical line, all squares between them are in  $B$ . Recently, it was proved in [3] that  $\alpha(B) = \theta(B)$  holds for any vertically convex board  $B$ . Considering Chung's counterexample in Figure 1, this is the most general possible version of the conjecture in some sense.

As it is written in [1], P. Erdős asked if  $\theta/\alpha$  is bounded and the answer is not known. Chung's example has  $\theta/\alpha = 8/7$ . The most that Chaiken et al. achieved for  $\theta/\alpha$  is  $21/17 - \varepsilon$  for any sufficiently small  $\varepsilon$ . Here we prove the following

**Theorem.**  $\theta(B) \leq 2\alpha(B) - 1$  for any simply connected board  $B$ .

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**Proof.** Let  $B$  be a simply connected board. Consider the intersections of the (closed) board with the horizontal lines. Every intersection consists of some disjoint intervals. Take the intervals containing a segment of the boundary of  $B$ . Let  $I_1, I_2, \dots, I_m$  denote these (maximal) intervals. Now we prove

$$(1) \quad \alpha(B) \cong m/2$$

and

$$(2) \quad \theta(B) \leq m-1$$

which yield the statement of Theorem.

First we prove that  $\alpha(B) \cong m/2$ . By the definition of the intervals  $I_1, I_2, \dots, I_m$ , there exists a unit square  $S_i$  in  $B$  such that the boundary of  $S_i$  and the intersection of  $I_i$  and the boundary of  $B$  have a common segment for  $i=1, 2, \dots, m$ . These unit squares  $S_i$  are either above or under the intervals  $I_i$ , so without loss of generality, we may suppose that for  $k \cong m/2$  indices, e.g. for  $i=1, 2, \dots, k$ ,  $S_i$  is above  $I_i$ . Now we prove that the unit squares  $S_1, S_2, \dots, S_k$  constitute an antirectangle.

If  $S_i$  and  $S_j$  are on different levels (i.e. in different rows), e.g.  $S_i$  is on a higher level than  $S_j$  then  $S_i$  and  $S_j$  cannot belong to the same rectangle because such a rectangle would have to contain the lower neighbouring unit square of  $S_i$  that does not belong to  $B$  by the definition of  $S_i$ .

If  $S_i$  and  $S_j$  ( $i \neq j$ ) are on the same level then the segment between the lower sides of the unit squares  $S_i$  and  $S_j$  contains a point  $P$  not belonging to  $B$  because the lower sides of  $S_i$  and  $S_j$  are segments of  $I_i$  and  $I_j$ , respectively and  $I_i$  and  $I_j$  are maximal segments belonging to  $B$  by definition. Then the unit square containing  $P$  in the row of  $S_i$  and  $S_j$  does not belong to  $B$  and so  $S_i$  and  $S_j$  cannot belong to the same rectangle. This completes the proof of inequality (1).

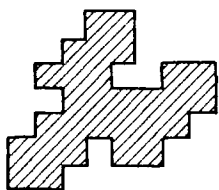


Fig. 1

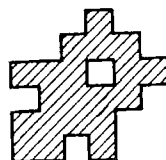


Fig. 2

Now we prove that  $\theta(B) \cong m-1$ . Let us consider the intervals  $I_1, I_2, \dots, I_m$ . Let us connect two intervals by a vertical segment if we can do it without crossing any other interval and without leaving the board  $B$ . Thus we obtain an arrangement  $A$  of vertical and horizontal segments in  $B$  without crossing. We prove that this arrangement does not contain any circuit (polygon line). Suppose that it is not the case. The angles of  $P$  are of 90 or 270 degrees. Suppose that there is an angle of 270 degrees. The horizontal arm of this angle is a subsegment  $S_0$  of a horizontal segment  $S$  of the arrangement  $A$ . The interior of  $P$  belongs to the board  $B$  by the simple connectivity of  $B$ , and so the subsegment  $S - S_0$  intersects the boundary of  $P$  by the maximality of  $S$ , contradicting the fact that the segments of the arrangement  $A$  do not intersect each other. Thus the angles of  $P$  are of 90 degrees,  $P$  is a rectangle and the horizontal sides of  $P$  are connected by at least two vertical segments, a contradiction.

Then consider the graph  $G$  with vertex-set  $V(G) = \{1, 2, \dots, m\}$  and edge-set  $E(G) = \{ij: I_i \text{ and } I_j \text{ are connected by a vertical segment}\}$ . Now  $G$  does not contain any circuit because it would correspond to a polygon  $P$  in the arrangement above. Thus the number  $e = |E(G)|$  of the connecting vertical segments is at most  $m-1$ . Let  $T_1, T_2, \dots, T_e$  denote these segments. Consider the maximal rectangles  $R_i$  containing  $S_i$  such that  $R_i$  is contained in  $B$  and that the orthogonal projection of  $R_i$  on the line of  $S_i$  is  $S_i$  for  $i=1, 2, \dots, e$ .

We prove that the rectangles  $R_1, R_2, \dots, R_e$  constitute a cover of  $B$ . Let  $X$  be an arbitrary point of  $B$ . Consider the vertical line  $L$  through  $X$ . Let  $I_a$  and  $I_u$  denote the intervals intersected by  $L$  first above and under  $X$ , respectively. Then  $I_a$  and  $I_u$  are connected by a vertical segment  $T_i \in \{T_1, \dots, T_e\}$  but they could have been connected by a segment  $T$  of  $L$ . Then  $T_i, T$  and the segments of  $I_a$  and  $I_u$  between  $T_i$  and  $T$  are in  $B$ , so by the simple connectivity of  $B$ , the whole rectangle bounded by these four segments is contained by  $B$ . But then  $R_i$  covers  $X$  and so we have proved that the rectangles  $R_1, \dots, R_e$  constitute a cover of  $B$ .

So we have proved that  $\theta(B) \leq e \leq m-1$  and the proof of the Theorem is complete. ■

### References

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